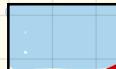


LECTURE 13

WEDNESDAY OCTOBER 23

Solving a Problem Recursively

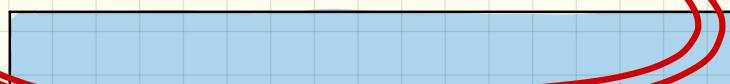
Given a **small** problem:



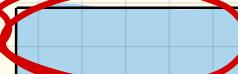
Solve it **directly**:



Given a **big** problem:



Divide it into **smaller** problems:



Assume solutions to **smaller** problems:



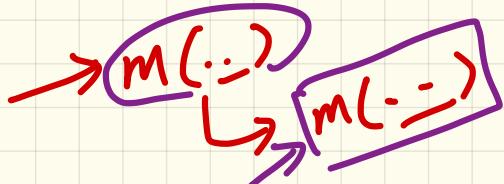
Combine solutions to **smaller** problems:



```
m(i) {  
    if(i == ...) /* base case: do something directly */ }  
    else {  
        m(j); /* recursive call with strictly smaller value */  
    }  
}
```

Tracing Recursion via a Stack

- When a method is called, it is **activated** (and becomes **active**) and **pushed** onto the stack.
- When the body of a method makes a (helper) method call, that (helper) method is **activated** (and becomes **active**) and **pushed** onto the stack.
 - ⇒ The stack contains activation records of all **active** methods.
 - **Top** of stack denotes the **current point of execution**.
 - Remaining parts of stack are (temporarily) **suspended**.
- When entire body of a method is executed, stack is **popped**.
 - ⇒ The **current point of execution** is returned to the new **top** of stack (which was **suspended** and just became **active**).
- Execution terminates when the stack becomes **empty**.



Runtime Stack

Problem

$$n! = \begin{cases} n=1 & 0 \\ n>1 \end{cases}$$

\$n\$ is size of original prob.
 n * $(n-1)!$ is size of strictly smaller prob.

→ size

4! = 4 × 3!

\times 3 × 2 × 1

3!

solution to a strictly smaller problem.

Recursive Solution: factorial

✓

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n \cdot (n - 1)! & \text{if } n \geq 1 \end{cases}$$

$$0! = 1$$

```

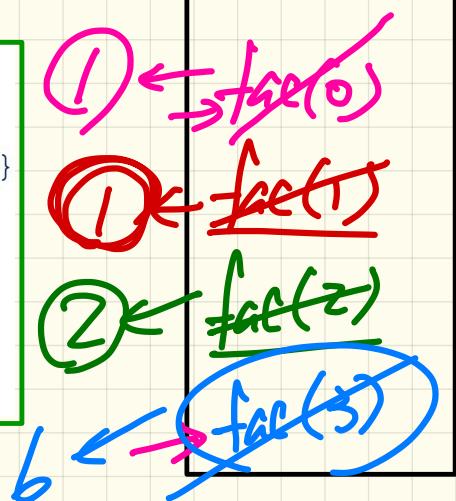
int factorial (int n) {
    int result;
    if (n == 0) { /* base case */ result = 1; }
    else { /* recursive case */
        result = n * factorial (n - 1);
    }
    return result;
}
    
```

3 2 1

Annotations:

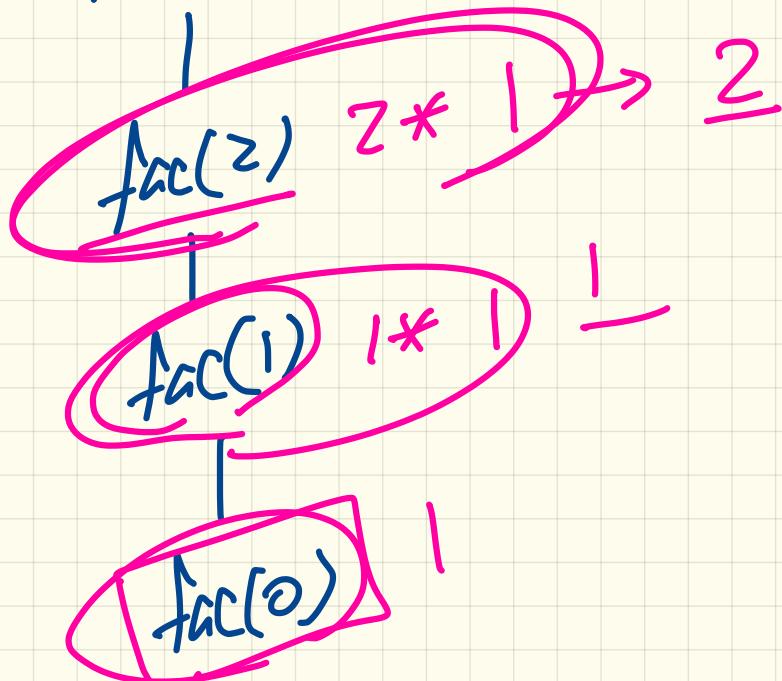
- Yellow boxes highlight `factorial`, `int`, and `n`.
- Handwritten numbers 3, 2, 1 are written above the code.
- Handwritten annotations show the recursive call `factorial(n - 1)` being multiplied by `n` to get the final result.
- Handwritten arrows point from the recursive call to the stack diagram.

Example: factorial(3)



$\rightarrow \text{fac}(3) \rightarrow * \text{ fac}(2)$

6.



Common Errors of Recursion (1)

```
int factorial (int n) {  
    return n * factorial (n - 1);  
}
```

fac(3)

↓
fac(-2)
fac(-1)
fac(0)
fac(1)
fac(2)
fac(3)



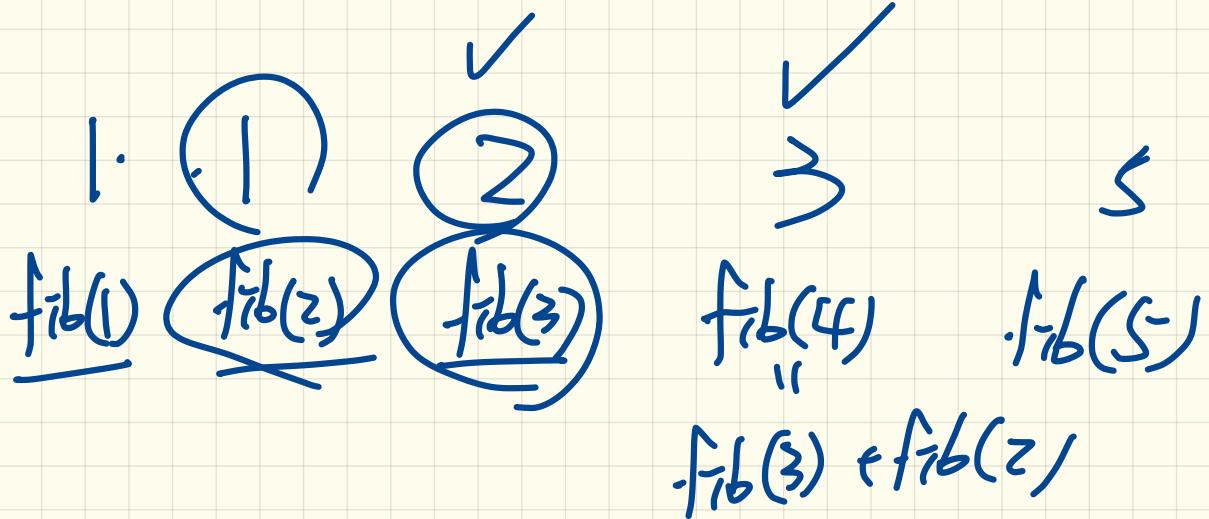
missing base
no termination

Common Errors of Recursion (2)

```
int factorial (int n) {  
    if (n == 0) /* base case */ return 1;  
    else { /* recursive case */ return n * factorial (n); }  
}
```

fact(3)

fact(3)
fact(3)
fact(3)
fact(3)



Recursive Solution: Fibonacci Number

$$F_n = \begin{cases} 1 & \text{if } n = 1 \\ 1 & \text{if } n = 2 \\ F_{n-1} + F_{n-2} & \text{if } n > 2 \end{cases}$$

combine.

↓ original problem

solution to another S.S. problem

```

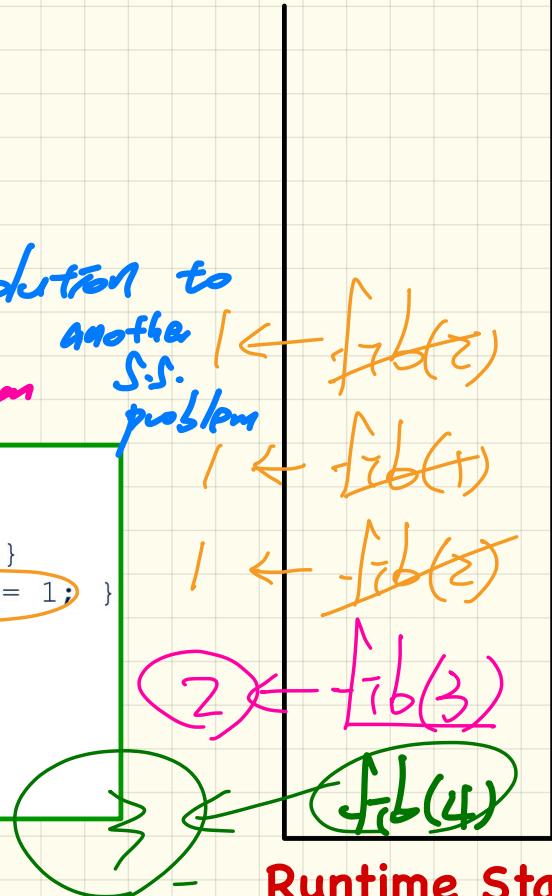
int fib (int n) {
    int result;
    if(n == 1) { /* base case */ result = 1; }
    else if(n == 2) { /* base case */ result = 1; }
    else { /* recursive case */
        result = fib(n - 1) + fib(n - 2);
    }
    return result;
}

```

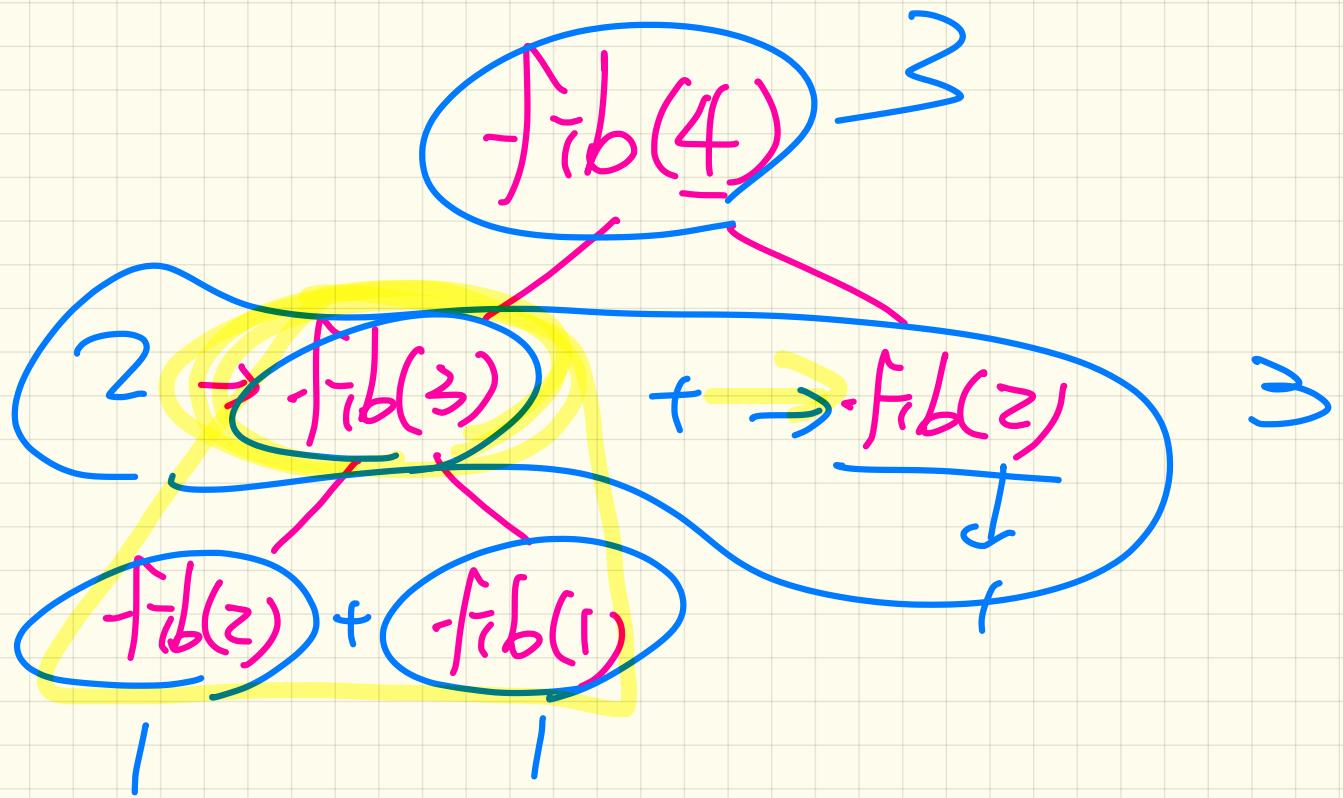
2 ← fib(3) + fib(2)

fib(2) + fib(1)

Example: fib(4)



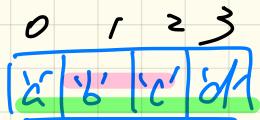
Runtime Stack

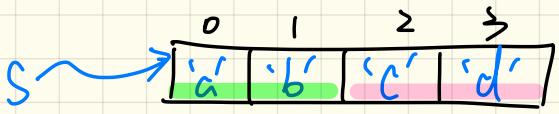


Use of String

O, -1

```
public class StringTester {  
    public static void main(String[] args) {  
        → String s = "abcd";  
        System.out.println(s.isEmpty()); /* false */  
        /* Characters in index range [0, 0) */  
        String t0 = s.substring(0, 0); → "" inclusive [0, 0)  
        System.out.println(t0); /* "" */  
        /* Characters in index range [0, 4) */  
        String t1 = s.substring(0, 4); → get from S[0] ~ S[3] exclusive  
        System.out.println(t1); /* "abcd" */  
        /* Characters in index range [1, 3) */  
        String t2 = s.substring(1, 3); → get from S[1] ~ S[2]  
        System.out.println(t2); /* "bc" */  
        String t3 = s.substring(0, 2) + s.substring(2, 4);  
        System.out.println(s.equals(t3)); /* true */  
        for(int i = 0; i < s.length(); i++) {  
            System.out.print(s.charAt(i));  
        }  
        System.out.println();  
    }  
}
```





String
(

s

$s.\text{substring}(0, 2)$

i is a valid
index

$+ s.\text{substring}(2, s.\text{length}())$

$s.\text{substring}(0, i)$

+

$s.\text{substring}(i, s.\text{length}())$

||

s

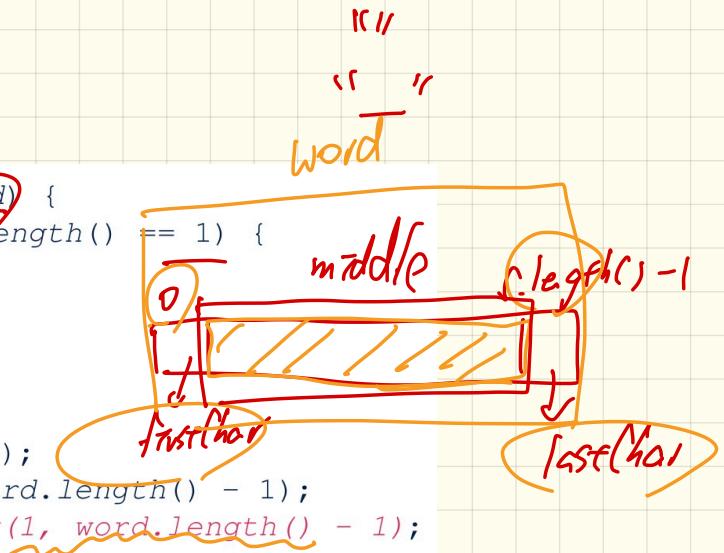
~~r~~ a c e c a ~~r~~

 |
 |
~~a~~ c
 |
 |
~~a~~ c a

The diagram illustrates the string "racecar" with various characters highlighted. The first character 'r' is crossed out with a red line. The string is enclosed in a pink rectangular border. Inside, the characters are highlighted: 'a' (green), 'c' (orange), 'e' (orange), 'c' (orange), 'a' (green). A red arrow points to the second 'e'. Below the string, the characters are mapped to their positions: the first 'a' is green, the second 'c' is orange, the third 'c' is orange, and the last 'a' is green. Orange lines connect the second 'c' and the third 'c' to their respective positions below, while green lines connect the first 'a' and the last 'a'.

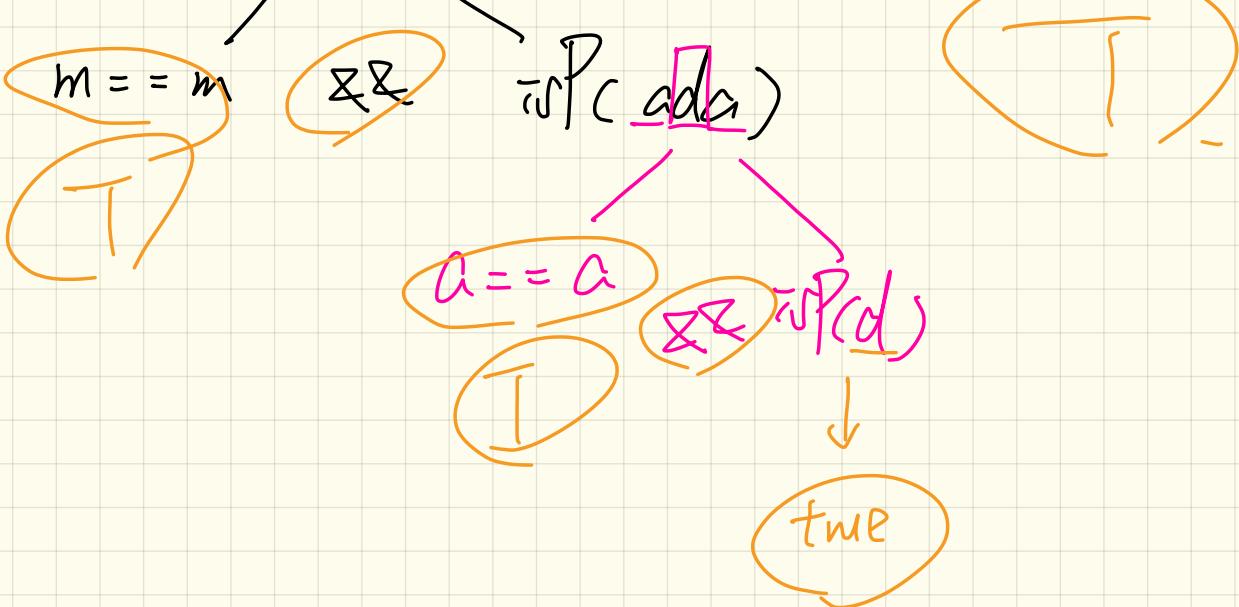
Problem: Palindrome

```
boolean isPalindrome (String word) {  
    if (word.length() == 0 || word.length() == 1) {  
        /* base case */  
        return true;  
    }  
    else {  
        /* recursive case */  
        char firstChar = word.charAt(0);  
        char lastChar = word.charAt(word.length() - 1);  
        String middle = word.substring(1, word.length() - 1);  
        return  
            firstChar == lastChar  
            /* See the API of java.lang.String.substring. */  
            && isPalindrome (middle);  
    }  
}
```



middle VS .
 word
 <

$\text{isP}(\underline{\text{m}}\underline{\text{a}}\underline{\text{d}}\underline{\text{a}}\underline{\text{m}})$



$$\rightarrow P(\underline{abc}\underline{a})$$

$G = \{G\}$ ~~QQ~~ 75 P(bc)

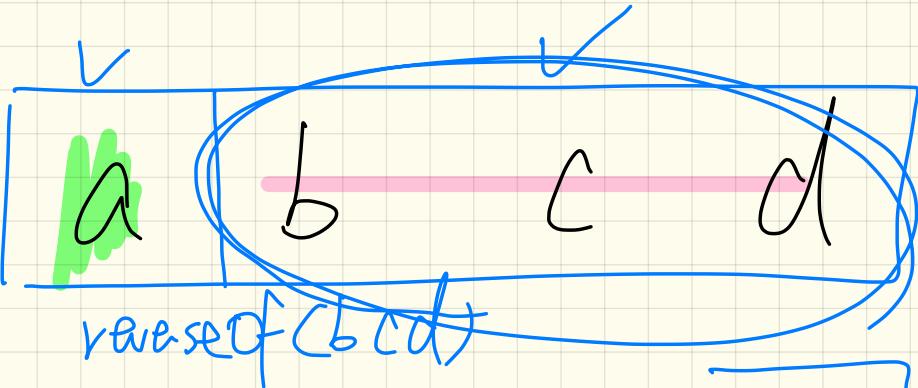
E

$b = -c$ $\Rightarrow P(c, \dots)$

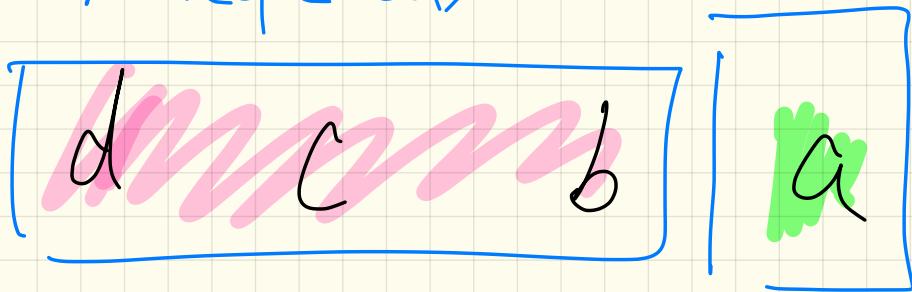
1

II

input →



output →



reverseOf (lefgh)

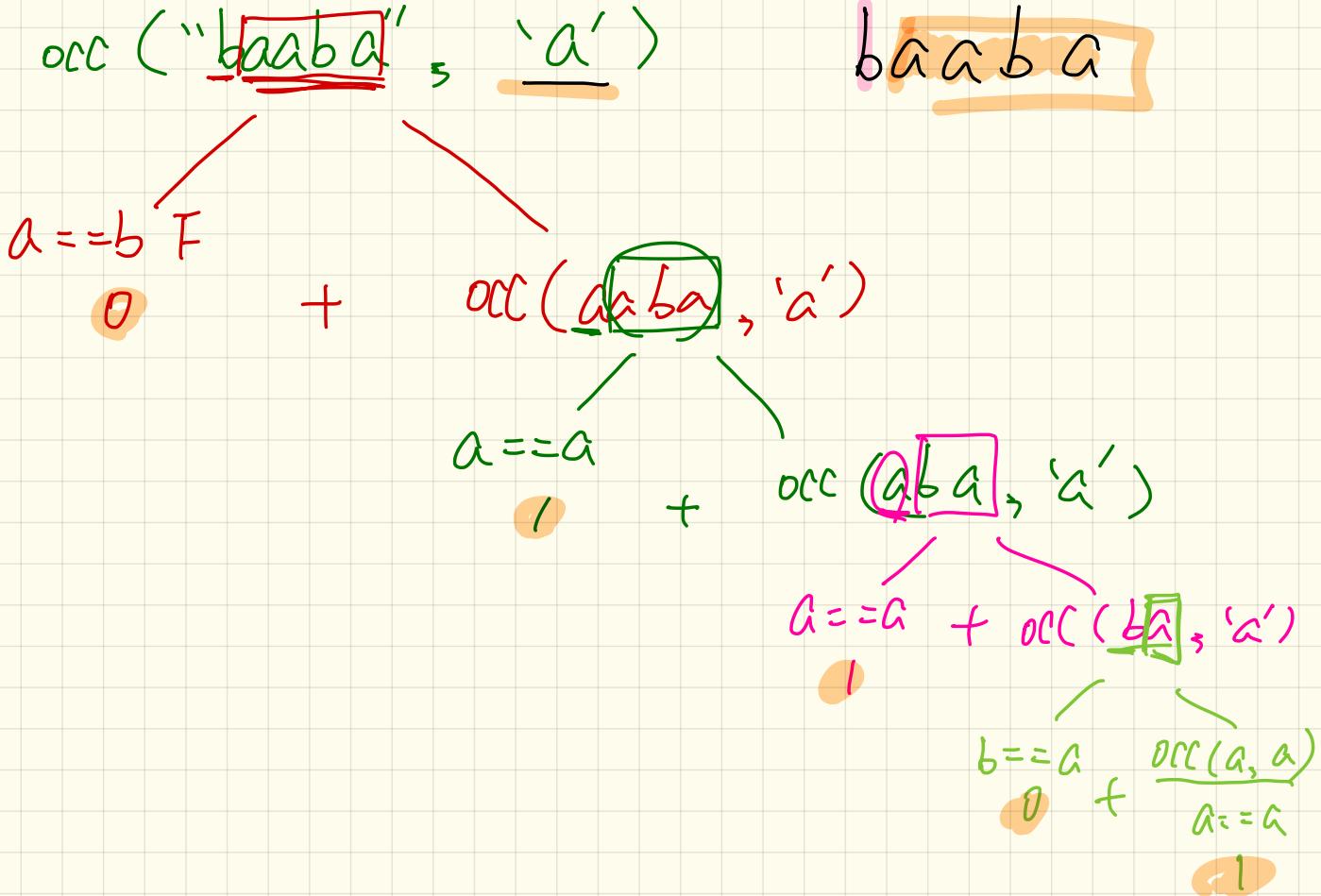
reverseOf (fgh) + e

reverseOf (gh) + f

↓
reverseOf (h)
h +

g

h g f e

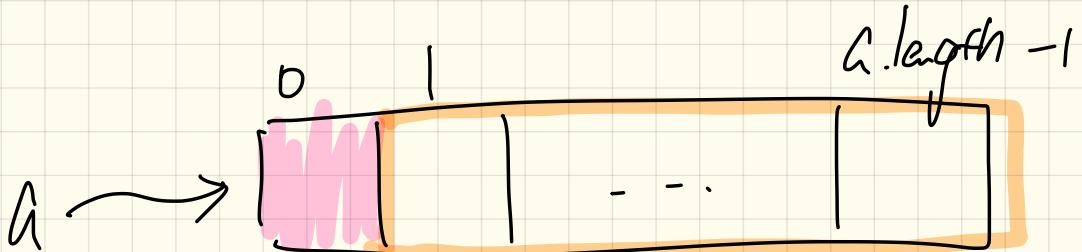


Problem: Reverse of a String

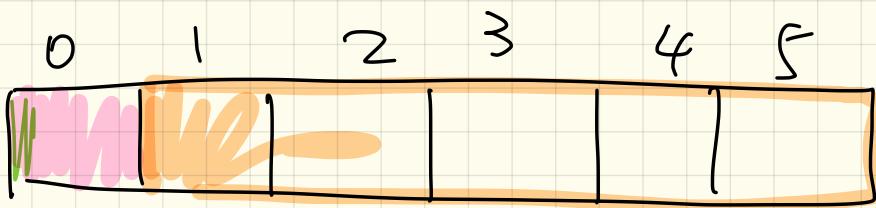
```
String reverseOf (String s) {  
    if(s.isEmpty()) { /* base case 1 */  
        return "";  
    }  
    else if(s.length() == 1) { /* base case 2 */  
        return s;  
    }  
    else { /* recursive case */  
        String tail = s.substring(1, s.length());  
        String reverseOfTail = reverseOf(tail);  
        char head = s.charAt(0);  
        return reverseOfTail + head;  
    }  
}
```

Problem: Number of Occurrences

```
int occurrencesOf (String s, char c) {  
    if (s.isEmpty()) {  
        /* Base Case */  
        return 0;  
    }  
    else {  
        /* Recursive Case */  
        char head = s.charAt(0);  
        String tail = s.substring(1, s.length());  
        if (head == c) {  
            return 1 + occurrencesOf (tail, c);  
        }  
        else {  
            return 0 + occurrencesOf (tail, c);  
        }  
    }  
}
```



0 1 5
5



2 5

3 5

4 5
(5-5) →